

# Logic Programming

## *Transitions*

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# Queries

True or False questions:

e.g. *Is Art the parent of Bob?*

Fill-in-the-blanks questions:

e.g. *Art is the parent of \_\_\_\_\_?*

e.g. *\_\_\_\_\_ is the parent of Bob?*

e.g. *\_\_\_\_\_ is the parent of \_\_\_\_\_?*

Compound questions:

e.g. *Is Art the parent of Bob **or** the parent of Bud?*

e.g. *\_\_\_\_\_ has sons **and** no daughters?*

# Values

Numbers and Strings:

e.g.  $2 + 2$ ?

e.g. *Concatenation of "abc" and "def"?*

Aggregates:

e.g. *Number of children of Art?*

e.g. *Number of people with daughters?*

Combinations:

e.g. *maximum of  $2+3$  and  $2*3$ ?*

e.g. *number Art's sons + number of Art's daughters?*

# Transforms

## Additions

e.g. *Add the fact that Art is the parent of Bill.*

## Deletions

e.g. *Drop the fact that Art is the parent of Bill.*

## Conditional Changes:

e.g. *Add grandparent facts implied by parent facts.*

e.g. *Replace parent facts with corresponding child facts.*

# Transitions

# Transforms

## Additions

e.g. *Add the fact that Art is the parent of Bill.*

## Removals

e.g. *Drop the fact that Art is the parent of Bill.*

## Conditional Changes:

e.g. *Add grandparent facts implied by parent facts.*

e.g. *Replace parent facts with corresponding child facts.*

# Update Rule

$$p(a,b) \ \& \ \underbrace{\sim q(b)}_{\text{conditions}} \implies \sim p(a,b) \ \& \ \underbrace{p(b,a)}_{\text{effects}}$$

# Update Rule

$$p(X, Y) \ \& \ \underbrace{\sim q(Y)}_{\text{conditions}} \implies \sim p(X, Y) \ \& \ \underbrace{p(Y, X)}_{\text{effects}}$$

# Safety

A transition rule is **safe** if and only if every variable in every literal on the right hand side appears in a positive literal on the left hand side. Also, every variable in a negative literal on the left hand side appears in a prior positive literal.

## Safe Transition Rule

$$p(X) \ \& \ \sim q(X) \ ==> \ \sim p(X) \ \& \ q(X)$$

## Unsafe Transition Rules

$$p(X) \ \& \ \sim q(X) \ ==> \ \sim p(X) \ \& \ q(Y)$$

$$p(X) \ \& \ \sim q(Y) \ ==> \ \sim p(X) \ \& \ q(X)$$

# Updates

An **update** is a finite collection of update rules.

## Example

$$p(X) \ \& \ \sim q(X) \ ==> \ \sim p(X) \ \& \ q(X)$$

$$p(X) \ \& \ q(X) \ ==> \ \sim p(X) \ \& \ \sim q(X)$$

# Instances

An **instance of an update rule** is a rule in which all variables have been consistently replaced by ground terms.

## Rule

$$p(X) \ \& \ \sim q(X) \ ==> \ \sim p(X) \ \& \ q(X)$$

## Herbrand Universe

$$\{a, b, c, d\}$$

## Instances

$$p(a) \ \& \ \sim q(a) \ ==> \ \sim p(a) \ \& \ q(a)$$

$$p(b) \ \& \ \sim q(b) \ ==> \ \sim p(b) \ \& \ q(b)$$

$$p(c) \ \& \ \sim q(c) \ ==> \ \sim p(c) \ \& \ q(c)$$

$$p(d) \ \& \ \sim q(d) \ ==> \ \sim p(d) \ \& \ q(d)$$

# Active and Inactive Rule Instances

A dataset  $\Delta$ , we say that update rule instance  $r$  is **active** if and only if  $\Delta$  contains all of the  $r$ 's positive conditions and none of its negative conditions. Otherwise, it is **inactive**.

**Rule:**  $p(X) \ \& \ \sim q(X) \implies \sim p(X) \ \& \ q(X)$

**Dataset:**  $\{p(b), p(c), p(d), q(d)\}$

**Instances** (active black, inactive grey):

$p(a) \ \& \ \sim q(a) \implies \sim p(a) \ \& \ q(a)$

$p(b) \ \& \ \sim q(b) \implies \sim p(b) \ \& \ q(b)$

$p(c) \ \& \ \sim q(c) \implies \sim p(c) \ \& \ q(c)$

$p(d) \ \& \ \sim q(d) \implies \sim p(d) \ \& \ q(d)$

# Positive and Negative Updates

The **positive update set**  $A(r, \Delta)$  for an update *rule*  $r$  on a dataset  $\Delta$  is set of all positive conclusions in some active instance of  $r$ . The **negative update set**  $D(r, \Delta)$  for an update *rule*  $r$  on a dataset  $\Delta$  is set of all negative conclusions in some active instance of  $r$ .

The **positive update set**  $A(\Omega, \Delta)$  for a *set of rules*  $\Omega$  on a dataset  $\Delta$  is the union of the positive updates for all of the rules in  $\Delta$ . The **negative update set**  $D(\Omega, \Delta)$  for a *set of rules*  $\Omega$  on a dataset  $\Delta$  is the union of the negative updates of all of the rules in  $\Delta$ .

# Update

The **result** of applying an update  $\Omega$  to a dataset  $\Delta$  (written  $u(\Omega, \Delta)$ ) is the set obtained by removing the negative updates and adding in the positive updates, i.e. the result is.

$$u(\Omega, \Delta) = \Delta - D(\Omega, \Delta) \cup A(\Omega, \Delta)$$

**Rule:**  $p(X) \ \& \ \sim q(X) \implies \sim p(X) \ \& \ q(X)$

**Dataset:**  $\{p(b), p(c), p(d), q(d)\}$

**Negative Update:**  $\{p(b), p(c)\}$

**Positive Update:**  $\{q(b), q(c)\}$

**Result:**  $\{p(d), q(b), q(c), q(d)\}$

# Interesting Example

## Update Rule:

$$p(X, a) \implies \sim p(X, a) \ \& \ p(a, X)$$

**Positive Update Set:**  $\{p(a, a)\}$

**Negative Update Set:**  $\{p(a, a)\}$

## Dataset:

$p(a, a)$

$p(a, b)$

$p(b, c)$

$p(c, c)$

$p(c, d)$

## New Dataset:

$p(a, a)$

$p(a, b)$

$p(b, c)$

$p(c, c)$

$p(c, d)$

# Interesting Example

## Update Rule:

$$p(X, a) \implies \sim p(X, a) \ \& \ p(a, X)$$

**Positive Update Set:**  $\{p(a, a)\}$

**Negative Update Set:**  $\{p(a, a)\}$

## Dataset:

$p(a, a)$

$p(a, b)$

$p(b, c)$

$p(c, c)$

$p(c, d)$

## New Dataset:

$p(a, a)$

$p(a, b)$

$p(b, c)$

$p(c, c)$

$p(c, d)$

# Production Systems

A **production system** is a set of condition-action rules. On each step in the execution of a production system, an active rule instance is chosen and the actions are performed. The cycle then repeats on the new state.

```
if p(X), then del p(X) and add p(X+10)
```

Before: {p(10), p(20), p(30)}

Step 1: {p(20), p(30)}

Step 2: {p(30)}

# Updates

Updates differ from production systems in that all active update rules “fire” at the same time. (1) All updates are computed *before* any changes are made, and (2) all changes are made simultaneously.

$p(X) \ \& \ \text{evaluate}(\text{plus}(X,10),Y) \implies \sim p(X) \ \& \ p(Y)$

Before:  $\{p(10), p(20), p(30)\}$

Step 1:  $\{p(20), p(30), p(40)\}$

Step 2:  $\{p(30), p(40), p(50)\}$

# The Game of Life



# Rules of the Game

- (1) Any *live* cell with *two or three* live neighbors lives on to the next generation.
- (2) Any *live* cell with *fewer than two* live neighbors dies (as if caused by underpopulation).
- (3) Any *live* cell with *more than three* live neighbors dies (as if by overpopulation).
- (4) Any *dead* cell with *exactly three* live neighbors becomes a live cell (as if by reproduction).

# Vocabulary

Symbols:  $c_{11}$ ,  $c_{12}$ , ...

Unary Predicates:

`on` - cell is live

`cell` - true of cells

Binary Predicates:

`neighbor` - cells are neighbors

# Social Starvation

Any *live* cell with *fewer than two* live neighbors dies.

`on(Y) & countofall(X,neighbor(X,Y)&on(X),0)`  
`==> ~on(Y)`

`on(Y) & countofall(X,neighbor(X,Y)&on(X),1)`  
`==> ~on(Y)`

# Overcrowding

Any *live* cell with *more than three* live neighbors dies.

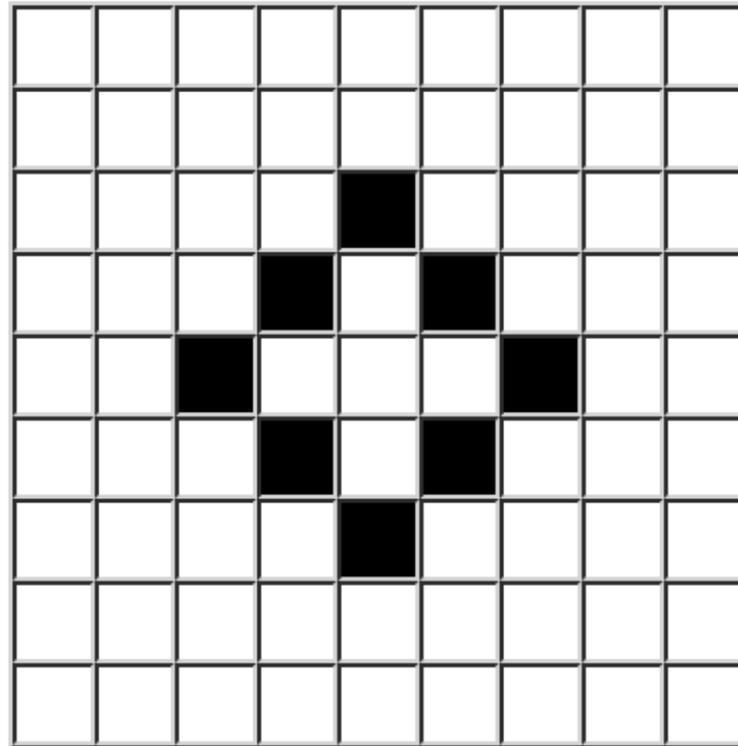
`on(Y) & countofall(X,neighbor(X,Y)&on(X),N) & leq(4,N)`  
`==> ~on(Y)`

# Reproduction

Any *dead* cell with *exactly three* live neighbors becomes live.

```
cell(Y) & ~on(Y) & countofall(X,neighbor(X,Y)&on(X),3)  
==> on(Y)
```

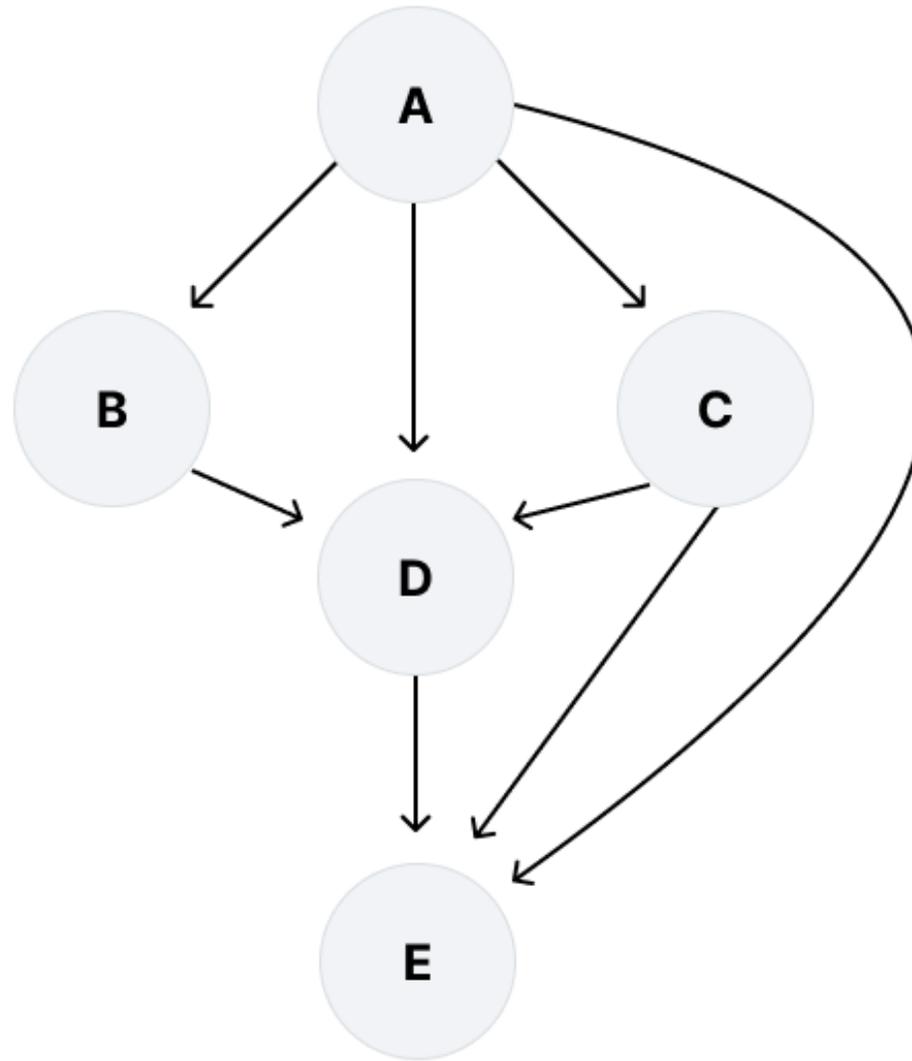
# Example



Demonstration

# Directed Graphs

# Example



# Copying

## Dataset:

$\{\text{edge}(a, b), \text{edge}(b, d), \text{edge}(b, e)\}$

## Operation Definition:

$\text{edge}(b, z) \implies \text{edge}(c, z)$

**Negative Updates:**  $\{\}$

**Positive Updates:**  $\{\text{edge}(c, d), \text{edge}(c, e)\}$

## Result:

$\{\text{edge}(a, b), \text{edge}(b, d), \text{edge}(b, e), \text{edge}(c, d), \text{edge}(c, e)\}$

# Reversing Arguments

## Dataset:

$\{\text{edge}(a, b), \text{edge}(b, d), \text{edge}(b, e)\}$

## Operation Definition:

$\text{edge}(X, Y) \implies \sim\text{edge}(X, Y) \ \& \ \text{edge}(Y, X)$

**Negative Updates:**  $\{\text{edge}(a, b), \text{edge}(b, d), \text{edge}(b, e)\}$

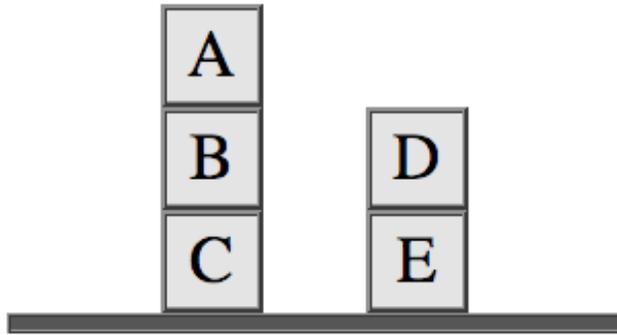
**Positive Updates:**  $\{\text{edge}(b, a), \text{edge}(d, b), \text{edge}(e, b)\}$

## Result:

$\{\text{edge}(b, a), \text{edge}(d, b), \text{edge}(e, b)\}$

# Blocks World

# Data



block(a)

block(b)

block(c)

block(d)

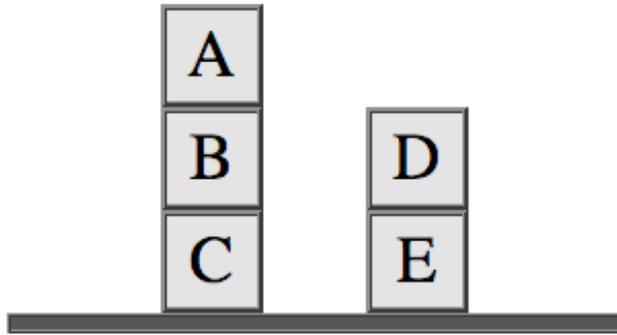
block(e)

on(a,b)

on(b,c)

on(d,e)

# Support and Clutter



```
block(a)
block(b)      on(a,b)
block(c)      on(b,c)
block(d)      on(d,e)
block(e)
```

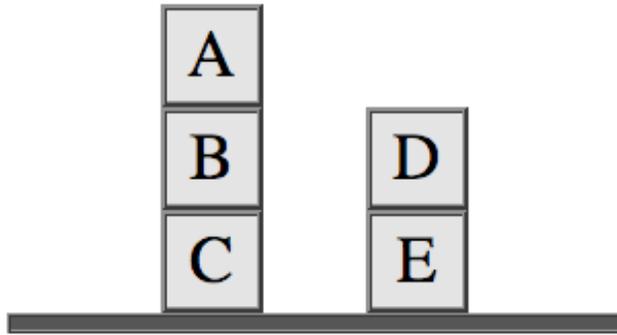
```
goal(X) :- on(X,Y)
```

```
goal(a)
goal(b)
goal(d)
```

```
goal(Y) :- on(X,Y)
```

```
goal(b)
goal(d)
goal(e)
```

# Materializing Relations

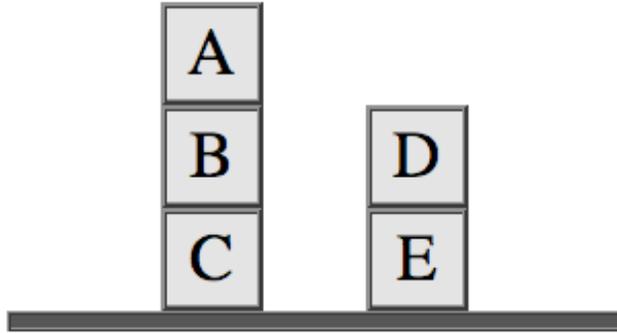


```
block(a)
block(b)      on(a,b)
block(c)      on(b,c)
block(d)      on(d,e)
block(e)
```

$\text{on}(X, Y) \implies \text{supported}(X) \ \& \ \text{cluttered}(Y)$

```
supported(a)
supported(b)
supported(d)
cluttered(b)
cluttered(c)
cluttered(e)
```

# Blocks World - above

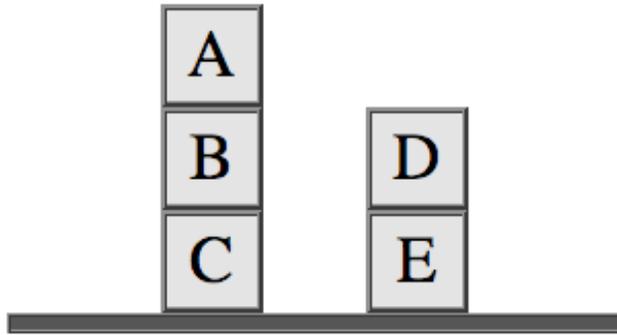


```
block(a)
block(b)      on(a,b)
block(c)      on(b,c)
block(d)      on(d,e)
block(e)
```

```
goal(X,Y) :- on(X,Y)
goal(X,Z) :- on(X,Y) & on(Y,Z)
goal(X,W) :- on(X,Y) & on(Y,Z) & on(Z,W)
...
```

```
goal(a,b)
goal(b,c)
goal(a,c)
goal(d,e)
```

# Blocks World - above



```
block(a)
block(b)      on(a,b)
block(c)      on(b,c)
block(d)      on(d,e)
block(e)
```

$\text{on}(X,Y) \implies \text{above}(X,Y)$

```
above(a,b)
above(b,c)
above(d,e)
```

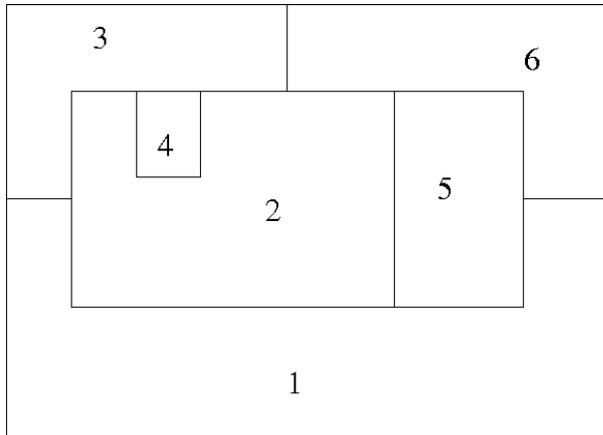
$\text{on}(X,Y) \ \& \ \text{above}(Y,Z) \implies \text{above}(X,Z)$

```
above(a,c)
```

*Repeat till fixpoint ("Bottom-up" computation of above.)*

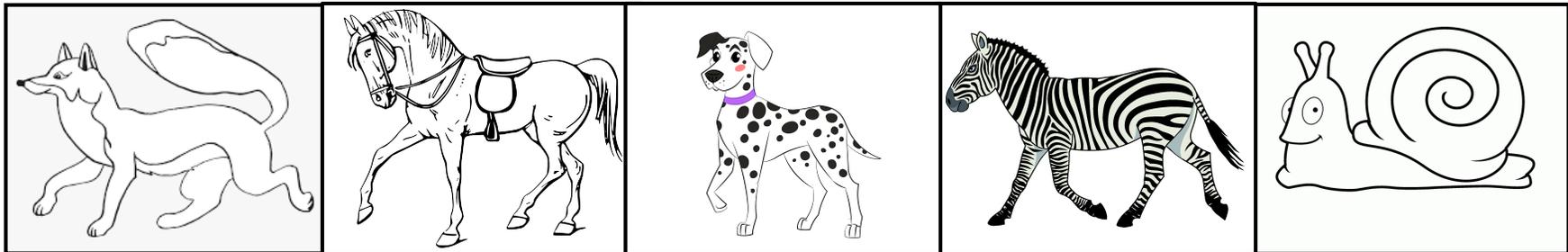
# Constraint Propagation

# Constraint Satisfaction Problems



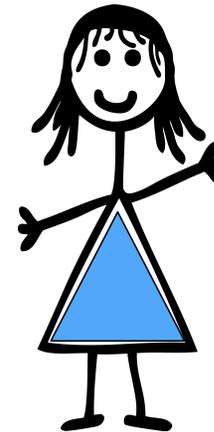
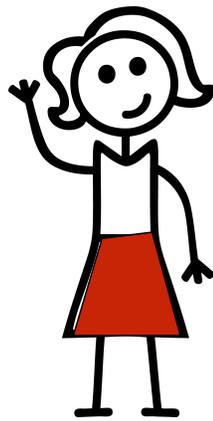
SEND  
+MORE  
-----  
MONEY

	6		1	4		5	
		8	3		5	6	
2							1
8			4	7			6
		6				3	
7			9	1			4
5							2
		7	2		6	9	
	4		5		8		7



# Problem Statement

Ruby Red, Willa White, Betty Blue are having lunch.  
One is wearing a red skirt, one white, one blue.

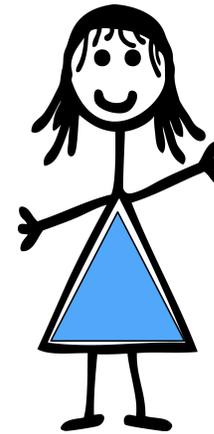
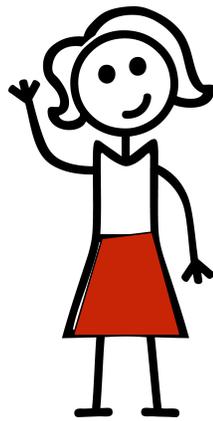


Betty to woman in white skirt: *Have you noticed that our skirts have colors different from our names?*

*Which woman is wearing which skirt?*

# Problem Statement

Ruby Red, Willa White, Betty Blue are having lunch.  
One is wearing a red skirt, one white, one blue.



Betty to woman in white skirt: *Have you noticed that our skirts have colors different from our names?*

*Why is this problem so easy to solve?*

# Vocabulary

Symbols: `red`, `white`, `blue`, `v`, `x`

Unary predicates Predicates:

`color(C)` is true if `C` is a color.

Ternary Predicates:

`c(P, C, v)` is true if person `P` is wearing color `C`.

`c(P, C, x)` is true if person `P` is not wearing color `C`.

Constraint Propagations Operations:

`c1` - No person is wearing same color as name.

`c2` - Betty Blue spoke to person in white.

`c3` - Every person has one color and vice versa.

# Colors

```
color(red)
```

```
color(white)
```

```
color(blue)
```

# First Given

No person is wearing same color as name.

`color(C) ==> c(C,C,x)`

# Second Given

Betty Blue spoke to person in white.

`c(blue,white,x)`

# Existence Constraints

Every person is wearing some color.

$$c(P, C1, x) \ \& \ c(P, C2, x) \ \& \ color(C3) \ \& \ mutex(C1, C2, C3) \\ \implies c(P, C3, v)$$

Every color is worn.

$$c(P1, C, x) \ \& \ c(P2, C, x) \ \& \ color(P3) \ \& \ mutex(P1, P2, P3) \\ \implies c(P3, C, v)$$

# Uniqueness Constraints

Nobody is wearing two colors.

$$c(P, C1, v) \ \& \ color(C2) \ \& \ distinct(C1, C2) \\ \implies c(P, C2, x)$$

No color is worn by two people.

$$c(P1, C, v) \ \& \ color(P2) \ \& \ distinct(P1, P2) \\ \implies c(P2, C, x)$$

# Initial State

Color

	r	w	b
Person			
r			
w			
b			

# First Given

No person is wearing same color as name.

`color(C) ==> c(C,C,x)`

	r	w	b
r			
w			
b			



	r	w	b
r	X		
w		X	
b			X

# Second Given

Betty Blue spoke to person in white.

`c(blue, white, x)`

	r	w	b
r	X		
w		X	
b			X



	r	w	b
r	X		
w		X	
b		X	X

# Existence Constraints

$c(P, C1, x) \ \& \ c(P, C2, x) \ \& \ color(C3) \ \& \ mutex(C1, C2, C3)$   
 $\implies c(P, C3, v)$

$c(P1, C, x) \ \& \ c(P2, C, x) \ \& \ human(P3) \ \& \ mutex(P1, P2, P3)$   
 $\implies c(P3, C, v)$

	r	w	b
r	X		
w		X	
b		X	X



	r	w	b
r	X	V	
w		X	
b	V	X	X

# Uniqueness Constraints

$c(P, C1, v) \ \& \ color(C2) \ \& \ distinct(C1, C2)$   
 $\implies c(P, C2, x)$

$c(P1, C, v) \ \& \ person(P2) \ \& \ distinct(P1, P2)$   
 $\implies c(P2, C, x)$

	r	w	b
r	X	V	
w		X	
b	V	X	X



	r	w	b
r	X	V	X
w	X	X	
b	V	X	X

# Existence Constraints

$c(P, C1, x) \ \& \ c(P, C2, x) \ \& \ color(C3) \ \& \ mutex(C1, C2, C3)$   
 $\implies c(P, C3, v)$

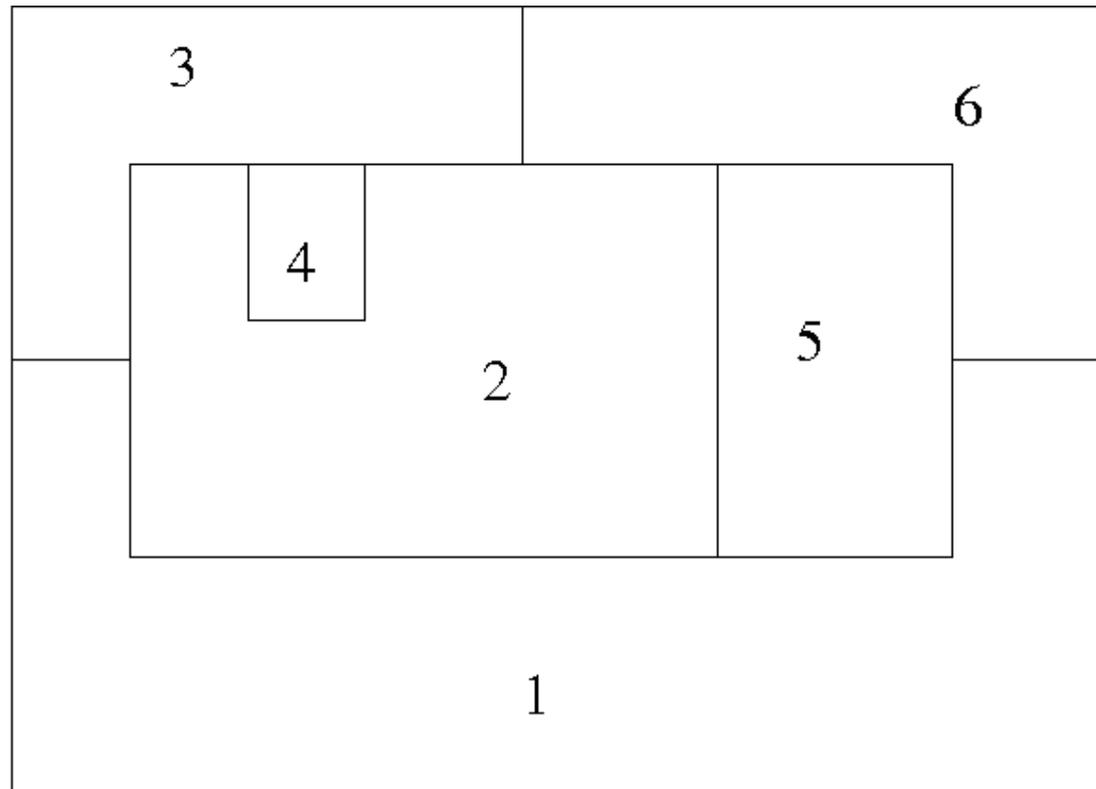
$c(P1, C, x) \ \& \ c(P2, C, x) \ \& \ human(P3) \ \& \ mutex(P1, P2, P3)$   
 $\implies c(P3, C, v)$

	r	w	b
r	X	V	X
w	X	X	
b	V	X	X



	r	w	b
r	X	V	X
w	X	X	V
b	V	X	X

# Map Coloring



# Sudoku

	6		1		4		5	
		8	3		5	6		
2								1
8			4		7			6
		6				3		
7			9		1			4
5								2
		7	2		6	9		
	4		5		8		7	



